Stochastic Sequential Screening

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Stochastic Allocation: Rationing and "Haggling"

- Rationing is common in practice.
- "Haggling" in marketplace interpreted as random allocation by Riley and Zeckhauser (1983) – is also common.
- Randomization helps relax the incentive constraints in the static environment (Myerson 1981, Bulow and Roberts 1989).
- Relatively little is known about the role of randomization in the dynamic setting.

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First-Order Approach and Deterministic Mechanism

- We know little because almost all papers in the (dynamic) mechanism design literature adopt the first-order approach.
 - form a relaxed problem by retaining only local downward IC.
 - impose some regularity conditions so that the solution to the relaxed problem has sufficient monotonicity to ensure global IC.
- Under these regularity conditions, optimal mechanisms are deterministic.

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Regularity and Monotonicity

- Regularity conditions are mild in static but not in dynamic setting.
- Static setting
 - static virtual surplus is increasing in agent's type,
 - ⇒ allocation monotone in type, which is also necessary for IC.
- Dynamic setting
 - dynamic virtual surplus is increasing in ex ante and ex post types,
 - ⇒ allocation is monotone is both ex ante and ex post types.

Problem: Monotonicity in ex ante type is not necessary; average monotonicity in ex ante type is.

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This Paper

Methodology

- we propose a method to deal with failure of regularity.
- we exploit necessity of average monotonicity of allocations.
- our method is valid whenever local IC implies global IC.

Results:

- show when and how randomization can increase revenue in a sequential screening setting.
- provide sufficient (and necessary) conditions for optimal mechanisms to be stochastic (\$\Rightarrow\$ new sufficient condition for optimal mechanisms to be deterministic).
- characterize optimal stochastic contracts.

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- Classical static mechanism design: Myerson (1981) and Riley and Zeckhauser (1983)
- Stochastic sequential screening: Courty and Li (2000), Krahmer and Kovac (2016), Krähmer and Strausz (2017), Battaglini and Lamba (2019), Krasikov and Lamba (2021)
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Model Setup

- Sequential screening with one seller, one buyer and one object:
 - buyer's ex ante type $\theta \in \{H, M, L\}$, with probability ϕ_{θ} , is a noisy signal of his true value (ex post type) $\omega \in [\underline{\omega}, \overline{\omega}]$;
 - first-order stochastic dominance: $F_H(\omega) \leq F_M(\omega) \leq F_L(\omega)$ for all ω ;
 - seller's reservation value $c \in (\omega, \overline{\omega})$;
 - direct mechanism: allocation rule $x_{\theta}(\omega)$, payment rule $t_{\theta}(\omega)$.

Remark: We need at least three ex ante types because deterministic contracts are proved to be optimal with binary ex ante types.

- Timing of the game:
 - **1** buyer observes θ ; seller proposes $\{x_{\theta}(\omega), t_{\theta}(\omega)\}$; buyer reports $\hat{\theta}$;
 - 2 buyer observes ω and reports $\hat{\omega}$; outcome $(x_{\hat{\theta}}(\hat{\omega}), t_{\hat{\theta}}(\hat{\omega}))$ is chosen

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Incentive and Participation Constraints

• Period-two IC $_{\theta}$: for all $\omega, \hat{\omega}$,

$$\omega x_{\theta}(\omega) - t_{\theta}(\omega) \ge \omega x_{\theta}(\hat{\omega}) - t_{\theta}(\hat{\omega}).$$

• Period-one $IC_{\theta\hat{\theta}}$: for all $\hat{\theta} \neq \theta$,

$$\int_{\underline{\omega}}^{\overline{\omega}} (\omega x_{\theta}(\omega) - t_{\theta}(\omega)) f_{\theta}(\omega) d\omega \ge \int_{\underline{\omega}}^{\overline{\omega}} (\omega x_{\hat{\theta}}(\omega) - t_{\hat{\theta}}(\omega)) f_{\theta}(\omega) d\omega.$$

• Period-one (interim) IR_{θ} :

$$\int_{\omega}^{\overline{\omega}} (\omega x_{\theta}(\omega) - t_{\theta}(\omega)) f_{\theta}(\omega) d\omega \ge 0.$$

Buyer's ex post payoff may fall below 0.

Seller's Problem

Seller chooses and commits to mechanism $\{x_{\theta}(\omega), t_{\theta}(\omega)\}$ to solve

$$\max_{(x_{\theta},t_{\theta})} \sum_{\theta=H.M.L} \phi_{\theta} \int_{\underline{\omega}}^{\overline{\omega}} (t_{\theta}(\omega) - cx_{\theta}(\omega)) f_{\theta}(\omega) d\omega$$

subject to IC_{θ} , $IC_{\theta\hat{\theta}}$ and IR_{θ} , for all $\theta, \hat{\theta} \in \{H, M, L\}$.

Simplified Problem

- As usual, no distortion at the top with $x_H(\omega) = \mathbb{1}\{\omega \geq c\}$.
- Choose weakly increasing $x_M(\omega)$ and $x_L(\omega)$ to maximize

$$\phi_M \int_{\underline{\omega}}^{\overline{\omega}} x_M(\omega) \delta_M(\omega) f_M(\omega) d\omega + \phi_L \int_{\underline{\omega}}^{\overline{\omega}} x_L(\omega) \delta_L(\omega) f_L(\omega) d\omega,$$

subject to an average monotonicity condition MON_{ML} :

$$\int_{\omega}^{\overline{\omega}} (x_M(\omega) - x_L(\omega)) (F_L(\omega) - F_M(\omega)) d\omega \ge 0,$$

where $\delta_{\theta}\left(\omega\right)$ are virtual surplus functions.

Optimal Deterministic Mechanism

- A mechanism with $x_{\theta}(\omega) = \mathbb{1}\{\omega \geq k_{\theta}\}$ is optimal iff $c \leq k_{M} \leq k_{L}$.
- Let (k_L, k_M) be the deterministic solution to the simplified problem:

$$\max_{k_M \leq k_L} \int_{k_M}^{\overline{\omega}} \phi_M \delta_M(\omega) f_M(\omega) d\omega + \int_{k_L}^{\overline{\omega}} \phi_L \delta_L(\omega) f_L(\omega) d\omega.$$

Then we have

$$(k_L,k_M) = \left\{ \begin{array}{ccc} (\hat{k}_L,\hat{k}_M) & \text{if} & \hat{k}_M \leq \hat{k}_L \\ (\hat{k},\hat{k}) & \text{with} & \hat{k} \in (\hat{k}_L,\hat{k}_M) & \text{if} & \hat{k}_M > \hat{k}_L \end{array} \right.$$

where \hat{k}_M , \hat{k}_L are unique maximizers of the two terms, respectively.

• Each (k_L, k_M) corresponds to an optimal deterministic mechanism.

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Toward Sufficient Condition for Randomization

• Failure of regularity $(\hat{k}_M > \hat{k}_L)$ is necessary for randomization, but is not sufficient, because it can be cheaper, in satisfying MON_{ML}

$$\int_{\omega}^{\overline{\omega}} \left(x_M \left(\omega \right) - x_L \left(\omega \right) \right) \left(F_L \left(\omega \right) - F_M \left(\omega \right) \right) d\omega \ge 0,$$

to set $k_L = k_M = \hat{k}$ than to randomize.

• If we can construct a stochastic mechanism that is more profitable than the deterministic solution \hat{k} , then optimal mechanisms must be stochastic.

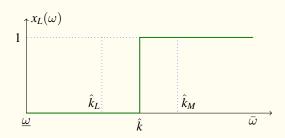
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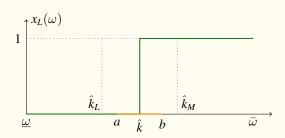
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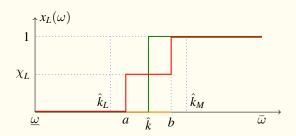
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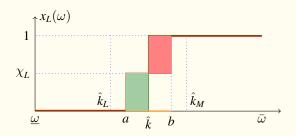
- Start from deterministic solution $k^M = k^L = \hat{k}$;
- Fix some interval $[a,b] \ni \hat{k}$ for type L (or type M);
- Set $x_L(\omega) = \chi_L \in (0,1)$ for all $\omega \in [a,b]$ to bind MON_{ML}: $\int_k^b (F_L(\omega) F_M(\omega)) d\omega = \chi_L \int_a^b (F_L(\omega) F_M(\omega)) d\omega;$
- Check if gain of (virtual) surplus over $[a, \hat{k}]$ exceeds loss over $[\hat{k}, b]$: $\int_{a}^{\hat{k}} \chi_{L} \phi_{L} \delta_{L}(\omega) f_{L}(\omega) d\omega \geq \int_{\hat{k}}^{b} (1 \chi_{L}) \phi_{L} \delta_{L}(\omega) f_{L}(\omega) d\omega.$



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Surplus-to-Rent Ratio

• The randomization in the construction improves revenue if

$$\frac{\int_a^{\hat{k}} \phi_L \delta_L(\omega) f_L(\omega) d\omega}{\int_a^{\hat{k}} (F_L(\omega) - F_M(\omega)) d\omega} \geq \frac{\int_{\hat{k}}^b \phi_L \delta_L(\omega) f_L(\omega) d\omega}{\int_{\hat{k}}^b (F_L(\omega) - F_M(\omega)) d\omega}.$$

• For any w' < w'', define the average surplus-to-rent ratio for θ as

$$R_{\theta}(w', w'') = \frac{\int_{w'}^{w''} \phi_{\theta} \delta_{\theta}(\omega) f_{\theta}(\omega) d\omega}{\int_{w'}^{w''} (F_{L}(\omega) - F_{M}(\omega)) d\omega}$$

• For any ω , define the point surplus-to-rent ratio ratio at ω as

$$r_{\theta}(\omega) = \frac{\phi_{\theta}\delta_{\theta}(\omega)f_{\theta}(\omega)}{F_{L}(\omega) - F_{M}(\omega)}$$

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Optimality of Deterministic Mechanisms

Theorem

Optimal mechanism is deterministic if either one of the following two conditions holds:

- $\hat{k}_M \leq \hat{k}_L$ (regularity condition);
- ② $\hat{k}_M > \hat{k}_L$, \hat{k} is interior, and the following condition holds for all θ :

$$\max_{\omega \le \hat{k}} R_{\theta}(\omega, \hat{k}) \le r_{\theta}(\hat{k}) \le \min_{\omega \ge \hat{k}} R_{\theta}(\hat{k}, \omega). \tag{1}$$

Remark: (i) Condition (1) is satisfied if $r_{\theta}(\omega)$ is increasing and is violated if $r_{\theta}(\omega)$ is decreasing. (ii) Condition (1) is necessary for randomization when regularity fails.

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Sufficient Conditions for Randomization

Define function $\alpha(\omega)$ as

$$F_{M}\left(\omega\right)=\left[1-\alpha\left(\omega\right)\right]F_{L}\left(\omega\right)+\alpha\left(\omega\right)F_{H}\left(\omega\right).$$

Theorem

Suppose $\hat{k}_M > \hat{k}_L$ and \hat{k} is interior. Any optimal mechanism is stochastic if either of the following two conditions holds: (i) α (ω) is weakly decreasing and

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(ii) $\alpha\left(\omega\right)$ is weakly increasing and

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Monotonicity of $\alpha\left(\omega\right)$ ensures that the stochastic solution to the simplified problem corresponds to an optimal mechanism.

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Optimal Randomization: Preliminary Observations

- At most one level of stochastic allocation for types M and L.
- Randomization for at most one type.
- If $r_{\theta}(\omega)$ is strictly increasing in some interval, then no randomization can happen in any subinterval for type θ ; if $r_{\theta}(\omega)$ is strictly decreasing on some interval, then allocation $x_{\theta}(\omega)$ cannot be deterministic with a threshold on the interval.

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Optimal Randomization: Characterization

Theorem

Suppose $r_{\theta}(\omega)$ is single-dipped for $\theta = L, M$, and appropriate boundary conditions are satisfied.

- (i) if $r_{\theta'}(\omega)$ is strictly increasing, then the simplified problem is solved by a stochastic allocation for θ and a deterministic allocation for θ' .
- (ii) if $r_{\theta'}(\omega)$ is strictly decreasing, then the simplified problem is solved by a deterministic allocation for θ and a stochastic allocation for θ' .

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Extension to Many Types

• Alignment condition (constant $\alpha(\omega)$): $\forall i \neq j, \forall i' \neq j'$, and $\forall \omega$,

$$\frac{f_{i}\left(\omega\right)-f_{j}\left(\omega\right)}{F_{i}\left(\omega\right)-F_{j}\left(\omega\right)}=\frac{f_{i'}\left(\omega\right)-f_{j'}\left(\omega\right)}{F_{i'}\left(\omega\right)-F_{j'}\left(\omega\right)}.$$

- Sufficient conditions for randomization can be found by the same construction, but they are no longer necessary.
- Ironing procedure for finding optimal stochastic mechanisms can be extended.

Conclusion

- We study randomization in a model of sequential screening.
- We provide necessary and sufficient conditions for randomization to be optimal.
- We characterize the optimal stochastic mechanism.