Miners' Reward elasticity and Stability of Competing Proof-of-Work Cryptocurrencies

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Summary

- Proof-of-Work cryptocurrencies (e.g., Bitcoin) hire miners to maintain the system by algorithmically setting the reward.
- □ Miners are freelance contributors and have strong discretion as to which cryptocurrency to contribute and how much they work for. Thus, the nature of miners' hash supply (≈ labor supply) is essential for the cryptocurrency's stability.
 □ Indeed, this paper (and our previous work, Noda, Okumura, and Hashimoto (2020)) shows that the combination of the difficulty adjustment algorithm (DAA, which controls miners' reward) and the value of the reward-elasticity of the hash supply is crucial for the cryptocurrency's stability.
 □ We develop a short-run supply-side model of the multicurrency mining market and estimate the hash supply elasticity of Bitcoin (BTC), Bitcoin Cash (BCH), and Bitcoin SV (BSV) by exploiting the discontinuity created by an event called halving.

Theoretical Prediction

- Preliminary results by Noda et al. (2020): In a single-currency model (i.e., we ignore BCH and BSV and focus only on BTC), the original DAA, CW-144, and ASERT asymptotically achieve the targeted average block time if and only if the (own) reward-elasticity of hash supply is smaller than 1, 144, and 575, respectively.
 If the elasticity is larger than one, the original DAA causes the overshoot of the winning rate. Too easy blocks (the winning rate is too low) and too difficult blocks (too high) arrive alternately, and the block time oscillates and diverges.
 Difficulty adjustment has two effects. (i) Change the speed of producing a
 - block by directly changing the difficulty of it. (ii) Change the hash rate by
- Bitcoin's DAA can stabilize the cryptocurrency only if the elasticity is low. The stability of Bitcoin hinges on external factors lowering the hash supply elasticity, such as the interaction with competing currencies (Bitcoin Cash, Bitcoin SV).
- By contrast, BCH and BSV are stable despite having a very elastic hash supply because they adopt efficient difficulty adjustment algorithms.
- By upgrading the difficulty adjustment algorithm, Bitcoin can prevent possible future crises before they happen.

Technical Background + Model

Cryptocurrency manages its transaction bistory using a ledger called blockchain.
 Blockchain is literally chain of blocks.

- We consider a multi-currency market. Each (crypto)currency is indexed by k.
- Block **height** is indexed by **t**.

□ **Miners** produce new blocks and append them to the blockchain.

 Image: Second state sta

changing the reward provided for miners. The original DAA completely dismisses (ii), and it does not work if (ii) is strong (i.e., hash supply is elastic).
 CW-144 and ASERT are convergent in virtually any environment.

□ In a multi-currency environment, the **cross elasticity** also matters.

- **Direct effect:** BTC's winning rate $\uparrow \Rightarrow$ BTC's hash rate \uparrow
- Indirect effect: BTC's winning rate $\uparrow \Rightarrow$ BCH's and BSV's hash rate $\downarrow \Rightarrow$ BCH and BSV increase their winning rate to maintain the block time = BCH's and BSV's reward $\uparrow \Rightarrow$ BTC's hash rate \downarrow
- The indirect effect attenuate the total elasticity
 - \rightarrow The existence of **BCH and BSV stabilizes BTC**'s block time.

Estimation + Simulation

U We approximate the hash supply function by a log-log linear function.

$$h(k,t;\alpha,\beta) = \bar{h}(t) \cdot \exp\left(\alpha_k + \sum_{k' \in [K]} \beta_{k',k} \log r(k',t)\right)$$

■ $\beta_{a,b}$ is currency *a*'s reward elasticity of currency *b*'s hash supply. □ We use MLE to estimate the parameters of (*α*, *β*).

\Box We identify β by looking at the data **before and after the third halving**.

Hash Supply (To)

- □ Many miners work on this task, and upon producing *t*-th block of currency *k*, the block creator receives a prize m(k, t) coins.
 - The value of m(k,t) for each t is prescheduled. It is halved every 210,000 blocks (this event is called halving). The last halving (third halving) occurred in 2020, and it reduces m from 12.5 to 6.25.
 - BCH (April 8, 2020) → BSV (April 10, 2020) → BTC (May 11, 2020)
- To prevent a miner from monopolizing the ledger, the system wants to randomly choose the next block creator.
- □ To this end, Proof-of-Work cryptocurrency requires miners to **draw lotteries**.
 - Draw a lottery = Computing a hash function once (counted as 1 hash).
 - **Hash rate** h(k, t) (hash/second) = **labor** input in a unit time.
 - The hash rate is not observable.
- □ The winning rate w(k, t) (= the probability of success per each lottery draw) is a policy variable, using a **difficulty adjustment algorithm** (**DAA**).
- Miner's expected reward from a unit hash computation: r(k,t) = w(k,t)m(k,t)e(k,t) (USD/hash), where e(k,t) is the exchange rate between the cryptocurrency k and USD.
 r(k,t) is publicly observable.
 We consider a short-run supply-side



| | | | | BTC | BCH | BSV | | |
|------------------------------|-----------|---------------|---------------------------------------|----------|-----------|-----------|--------------------|-----------------------------------|
| BTC | 52.879*** | Reward (From) | BTC | 0.626*** | -3.981*** | -3.186*** | - * ** _ *** | p < 0.05 p < 0.01 p < 0.001 |
| | (1.973) | | | (0.103) | (0.113) | (0.106) | | |
| BCH | 49.851*** | | BCH | -0.240* | 5.386*** | -1.540*** | | |
| | (1.995) | | | (0.095) | (0.127) | (0.093) | | |
| BSV | 47.764*** | | BSV | -0.223* | -1.219*** | 4.869*** | | |
| | (1.973) | | | (0.098) | (0.076) | (0.118) | | |
| Constant (α) | | | Own- and Cross-Elasticity (β) | | | | | |

- □ The hash supply is increasing in its own reward (diagonal elements) and decreasing in its rivals' reward (off-diagonal elements).
- BTC's hash supply is very inelastic (own elasticity < 1). Therefore, BTC has survived despite it has used the inefficient original DAA.</p>
- **BCH**'s and **BSV**'s own elasticities are much larger than 1. They were not to survive if the original DAA were maintained.



- **model** of miners. Miners' capital (facility for mining) is fixed, and miners decide how to **operate** dynamically.
- h(k,t) is a function of $(r(k,t))_{k\in K}$.

Club 256 Mining ASIC Can be used (exclusively) to mine

SHA-256 Miners BTC, BCH, or BSV

Electricity cost (variable)

□ Block time b(k, t) (= time needed for producing one block) approximately follows an **exponential distribution** with mean $\mathbb{E}[b(k, t)] = 1/w(k, t)h(k, t)$ (second). BTC, BCH, and BSV aim at achieving $\mathbb{E}[b(k, t)] = 600$ seconds.

Difficulty Adjustment Algorithm (DAA) selects a new winning rate w(k, t + 1) using past block time b(k, t') and winning rate w(k, t') (for $t' \le t$) as its inputs.

Multiple DAAs have been used and implemented. Original DAA,
 CW-144, and ASERT are the names of three different DAAs that are studied in this paper.



(a) BTC: actual

(b) BCH: actual

(c) BSV: actual

□ The simulation shows that if BTC faces a larger $\beta_{BTC,BTC}$, BTC's block time starts to **oscillate** and **diverge**. The threshold is around 1.5, which is substantially larger than 1.0 (thanks to the interaction with BCH and BSV).

